

# How to choose the number of taps in a DFE ?

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## Abstract

In this paper, we address the following problem. Given a fixed *total* number of feedforward and feedback filter taps in a decision feedback equalizer, what is the optimum number of taps for each filter and what is the optimum decision delay setting ? We propose a simple algorithm that reduces the solution search space from being 2-dimensional to 1-dimensional at a small performance loss. We apply our algorithm to indoor wireless channels.

## I Introduction

In an increasing number of applications, a minimum mean square error decision feedback equalizer (MMSE-DFE) is implemented in the receiver to mitigate intersymbol interference (ISI) on bandlimited dispersive channels. In most practical situations, computational and implementational complexity considerations often place a constraint on the maximum number of total (feedforward and feedback) filter taps that can be used. In addition, for short filters, the decision delay setting could affect performance significantly [1]. Therefore, it becomes of chief importance to understand the dependence of the finite-length MMSE-DFE performance on the number of its feedforward and feedback filters and on the decision delay in order to set them properly.

Currently-used techniques for setting these 3 parameters are either ad-hoc or computationally intense (require an exhaustive search over these 3 parameters and inversion of a matrix whose size is equal to the total number of feedforward and feedback taps for each step in the search). In this paper, we present a simple characterization of the MMSE-DFE's performance that results in a simple order-recursive algorithm. Furthermore, we reduce the exhaustive search space from a 2-dimensional to a 1-dimensional search for joint selection of the number of MMSE-DFE taps and the decision delay setting.

This paper is organized as follows. Section 2 presents closed-form expressions for computing the optimum DFE taps and evaluating its performance

given the number of feedforward and feedback taps and the decision delay. Section 3 proposes a simple 1-dimensional suboptimal search algorithm for setting these 3 parameters. Numerical results illustrating the near-optimum performance of the proposed algorithm are given in Section 4 and the paper is concluded in Section 5.

## II The Finite-Length MMSE-DFE

The channel impulse response (CIR) is assumed to be a finite impulse response (FIR) filter with memory  $\nu$ , i.e.,  $\mathbf{h} = [h_0 \ h_1 \ \dots \ h_\nu]^t$ , where  $(\cdot)^t$  denotes the transpose. Over a block of  $(N + 1)$  channel output symbols, denoted by  $\mathbf{y}_{i-N}^i$ , the input-output relationship can be cast in matrix form as follows (see also Equation (11) in the Appendix)

$$\mathbf{y}_{i-N}^i = \mathbf{H}\mathbf{x}_{i-N-\nu}^i + \mathbf{n}_{i-N}^i, \quad (1)$$

where  $\mathbf{H}$  is the  $(N + 1) \times (N + \nu + 1)$  Toeplitz <sup>1</sup> channel convolution matrix. The input and noise sequences (denoted by  $\mathbf{x}_{i-N-\nu}^i$  and  $\mathbf{n}_{i-N}^i$ , respectively) are assumed to be zero-mean, independent of each other, and have positive-definite (non-singular) auto-correlation matrices denoted by  $\mathbf{R}_{xx}$  and  $\mathbf{R}_{nn}$ , respectively. Using (1), the input-output cross-correlation matrix and the output auto-correlation matrix are given by

$$\mathbf{R}_{xy} \stackrel{def}{=} E[\mathbf{x}_{i-N-\nu}^i \mathbf{y}_{i-N}^{i*}] = \mathbf{R}_{xx} \mathbf{H}^* \quad (2)$$

$$\mathbf{R}_{yy} \stackrel{def}{=} E[\mathbf{y}_{i-N}^i \mathbf{y}_{i-N}^{i*}] = \mathbf{H} \mathbf{R}_{xx} \mathbf{H}^* + \mathbf{R}_{nn}, \quad (3)$$

where  $E[\cdot]$  and  $(\cdot)^*$  denote the expected value and the complex-conjugate transpose, respectively.

## A Optimum 2-Dimensional Search Algorithm

A block diagram of the MMSE-DFE is shown in Figure 3. The error symbol at time  $i$  is given by <sup>2</sup>  $e_i \stackrel{def}{=} \mathbf{y}_{i-N}^i - \hat{\mathbf{y}}_{i-N}^i$

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<sup>1</sup>The channel is assumed *time-invariant* over the transmission block.

<sup>2</sup>Assuming correct previous decisions, i.e.,  $\hat{x}_{i-\Delta} = x_{i-\Delta}$ .

$\hat{x}_{i-\Delta} - z_i = \tilde{\mathbf{b}}^* \mathbf{x}_{i-N-\nu}^i - \mathbf{w}^* \mathbf{y}_{i-N}^i$ , where  $\Delta$  is the decision delay,  $\tilde{\mathbf{b}}^* \stackrel{def}{=} [\mathbf{0}_{1 \times \Delta} \quad \mathbf{b}^* \quad \mathbf{0}_{1 \times s}]$ ,  $\mathbf{0}_{m \times n}$  is the all-zeros matrix with  $m$  rows and  $n$  columns, and  $s \stackrel{def}{=} N + \nu - \Delta - N_b$ . We assume that the feedforward filter consists of  $(N + 1)$  taps, i.e.,  $\mathbf{w} \stackrel{def}{=} [w_0 \quad \dots \quad w_N]^t$ , and the *strictly-causal* feedback filter  $\mathbf{e}_\Delta^* - \tilde{\mathbf{b}}^* = [\mathbf{0}_{1 \times (\Delta+1)} \quad -b_1 \quad \dots \quad -b_{N_b} \quad \mathbf{0}_{1 \times s}]$  consists of  $N_b$  non-zero taps, where  $\mathbf{b} \stackrel{def}{=} [1 \quad b_1 \quad \dots \quad b_{N_b}]^t$  and  $\mathbf{e}_\Delta^* = [\mathbf{0}_{1 \times \Delta} \quad 1 \quad \mathbf{0}_{1 \times (N+\nu-\Delta)}]$ .

Using the *Orthogonality Principle* of linear least-squares estimation, it can be shown that the mean square of the (assumed stationary) error sequence is equal to

$$\begin{aligned} MSE &\stackrel{def}{=} \mathbb{E}[|e_i|^2] \\ &= \tilde{\mathbf{b}}^* (\mathbf{R}_{xx} - \mathbf{R}_{xy} \mathbf{R}_{yy}^{-1} \mathbf{R}_{yx}) \tilde{\mathbf{b}} \\ &\stackrel{def}{=} \tilde{\mathbf{b}}^* \mathbf{R}_{x/y}^\perp \tilde{\mathbf{b}} \\ &= \mathbf{b}^* \mathbf{R}_\Delta \mathbf{b}, \end{aligned} \quad (4)$$

where  $\mathbf{R}_\Delta$  is an  $(N_b + 1) \times (N_b + 1)$  submatrix of  $\mathbf{R}_{x/y}^\perp$ .

It can be shown that the optimum feedback and feedforward filters that minimize (5) are given by

$$\mathbf{b}_{opt} = \frac{\mathbf{R}_\Delta^{-1} \begin{bmatrix} 1 \\ \mathbf{0}_{N_b \times 1} \end{bmatrix}}{\mathbf{R}_\Delta^{-1}(1, 1)} \quad (6)$$

$$\mathbf{w}_{opt} = \tilde{\mathbf{b}}_{opt}^* \mathbf{R}_{xy} \mathbf{R}_{yy}^{-1}, \quad (7)$$

where  $\mathbf{R}_\Delta^{-1}(1, 1)$  is the  $(1, 1)$  entry of  $\mathbf{R}_\Delta^{-1}$ . The resulting minimum mean square error (MMSE) is given by

$$MMSE(N, N_b, \Delta) = \frac{1}{\mathbf{R}_\Delta^{-1}(1, 1)}. \quad (8)$$

Therefore, the *unbiased* [2] decision-point SNR of the MMSE-DFE is equal to

$$SNR_{MMSE-DFE,U}(N, N_b, \Delta) = \frac{1}{MMSE(N, N_b, \Delta)} - 1. \quad (9)$$

Hence, for any choice of the triplet  $(N, N_b, \Delta)$ , Equations (6) and (7) are used to determine the optimum MMSE-DFE settings and Equation (9) is used to evaluate its performance. Under the implementation constraint  $N + N_b = N_{tot}$  (where  $N_{tot}$  is a given constant), determining the optimum triplet  $(N, N_b, \Delta)$  is computationally intense since it requires an exhaustive 2-dimensional<sup>3</sup> search and computation of the inverse  $\mathbf{R}_\Delta^{-1}$  for each step of the search. Next, we show how to reduce this computational complexity at a small performance loss.

<sup>3</sup>The linear constraint  $N + N_b = N_{tot}$  reduces the exhaustive search space for  $(N, N_b, \Delta)$  from being 3-dimensional to 2-dimensional.

## B Recursive Formula for MMSE

Using Kramer's rule for matrix inversion [3], we get the following alternative MMSE expression

$$MMSE(N, N_b, \Delta) = \frac{|\mathbf{R}_\Delta|}{|\tilde{\mathbf{R}}_\Delta^{1,1}|}, \quad (10)$$

where  $\tilde{\mathbf{R}}_\Delta^{i,j}$  is obtained from  $\mathbf{R}_\Delta$  by deleting its  $i^{th}$  row and  $j^{th}$  column and  $|\cdot|$  denotes the determinant. This compact characterization of the MMSE offers further insight into the effect of the triplet  $(N, N_b, \Delta)$  on the MMSE-DFE performance as illustrated in Figure 1 and described in the sequel :

- Given  $\mathbf{h}$ ,  $\mathbf{R}_{xx}$ , and  $\mathbf{R}_{nn}$ , the number of feedforward taps  $(N + 1)$  determines  $\mathbf{R}_{x/y}^\perp$  as defined in (4).
- The number of feedback taps  $N_b$  determines the *size* of  $\mathbf{R}_\Delta$ .
- The delay of the feedforward filter  $\Delta$  determines the *location* of the  $(1, 1)$  entry of  $\mathbf{R}_\Delta$ . Hence, changing the delay corresponds to *shifting*  $\mathbf{R}_\Delta$  along the main diagonal of  $\mathbf{R}_{x/y}^\perp$ .
- For a given  $N$ , optimizing  $N_b$  and  $\Delta$  is equivalent to locating the  $(N_b + 1)$ -dimensional matrix  $\mathbf{R}_\Delta$  whose upper-left corner is on the main diagonal of  $\mathbf{R}_{x/y}^\perp$  such that the determinant ratio<sup>4</sup> in (10) is minimized. Note that (10) can be computed using the well-known *order-recursive* formula for the determinant  $|\mathbf{R}_\Delta| = \sum_{j=1}^{N_b+1} (-1)^{j+1} \mathbf{R}_\Delta(1, j) |\tilde{\mathbf{R}}_\Delta^{1,j}|$ .

Finally, applying Kramer's rule to (6), we get the following alternative expression for the optimum feedback filter taps

$$\mathbf{b}_{opt} = \left[ 1 \quad \frac{|\tilde{\mathbf{R}}_\Delta^{1,2}|}{|\tilde{\mathbf{R}}_\Delta^{1,1}|} \quad \frac{|\tilde{\mathbf{R}}_\Delta^{1,3}|}{|\tilde{\mathbf{R}}_\Delta^{1,1}|} \quad \dots \quad \frac{|\tilde{\mathbf{R}}_\Delta^{1,(N_b+1)}|}{|\tilde{\mathbf{R}}_\Delta^{1,1}|} \right].$$

## III Suboptimum 1-Dimensional Search Algorithm

### Problem Statement

Given the CIR  $\mathbf{h}$  of memory  $\nu$  and a fixed total number of DFE taps  $N_{tot} \geq \nu$ <sup>5</sup>. Assume that the noise is white ( $\mathbf{R}_{nn} = \frac{1}{SNR} \mathbf{I}_{N+1}$  where  $SNR$  is the signal to noise ratio) and the input data is white ( $\mathbf{R}_{xx} = \mathbf{I}_{N+\nu+1}$ ). We want to solve the non-linear integer programming problem

$$\min_{N, N_b, \Delta} MMSE(N, N_b, \Delta) \text{ subject to } N + N_b = N_{tot},$$

<sup>4</sup>This ratio is also equal to the first diagonal element in an upper-diagonal-lower Cholesky factorization of  $\mathbf{R}_\Delta$ .

<sup>5</sup>When  $N_{tot} < \nu$ , the DFE performance is degraded significantly. Hence, this case is not useful in practice and is not considered.

where the *integers*  $N$ ,  $N_b$ , and  $\Delta$  satisfy

$$0 \leq N \leq N_{tot} - 1; 1 \leq N_b \leq N_{tot}; 0 \leq \Delta \leq N + \nu.$$

An exhaustive search would require  $\sum_{N=0}^{N_{tot}-1} (N + \nu + 1) = \frac{(N_{tot}+1)N_{tot}}{2} + N_{tot}\nu$  steps since there are  $(N + \nu + 1)$  possible delay values. We can further restrict the search space by noting that

- $1 \leq N_b \leq \nu$  (we prove in the Appendix that for  $N_b > \nu$  the last  $(N_b - \nu)$  feedback taps are equal to zero).
- $N \leq \Delta \leq N + \nu - N_b$  (since for  $N > \Delta$  the feedforward filter coefficients  $\mathbf{w}_{\Delta+1}^N$  are equal to zero. Moreover, we do not need  $\Delta > (N + \nu - N_b)$  because of the feedback filter post-cursor cancellation action).

Over these reduced ranges, the exhaustive search size is reduced <sup>6</sup> to  $\sum_{N_b=1}^{\nu} (\nu - N_b + 1) = \frac{\nu^2 + \nu}{2}$  steps but still has quadratic complexity. The following suboptimum algorithm reduces the search space to become 1-dimensional with  $\nu$  search steps.

#### Proposed Algorithm

**Input :** Total number of DFE tap  $N_{tot}$  and channel memory  $\nu$

1. For  $N_b = 1, 2, \dots, \nu$ .
2. Set  $N = N_{tot} - N_b$ .
3. Set  $\Delta = N + \nu - N_b = N_{tot} - 2N_b + \nu$ . In other words, *we assume that the feedback filter cancels all post-cursor ISI.*<sup>7</sup> Figure 2 illustrates this assumed relationship between the DFE parameters.
4. Choose the value of  $N_b$  that results in the minimum MMSE, as given in (10), i.e.

$$N_b = \underset{1 \leq N_b \leq \nu}{\operatorname{argmin}} MMSE(N = N_{tot} - N_b, N_b, \Delta = N_{tot} - 2N_b + \nu).$$

The corresponding values of  $N$  and  $\Delta$  are determined from Steps 2 and 3 above. Figure 4 depicts an example of the variation of  $SNR_{MMSE-DFE,U}$  versus  $N_b$  for a wireless indoor channel at input SNR of 15 dB with  $N_{tot} = 10$ . The existence of several local optima justifies the need for the 1-dimensional search in general. However, for large  $N_{tot}$  (compared to  $\nu$ ) no search is needed as detailed next.

#### Special Case

When  $N_{tot} \gg \nu$ , the optimum values for the triplet  $(N, N_b, \Delta)$  are  $N_b = \nu$ ,  $N = N_{tot} - \nu$ , and  $\Delta = N$ .<sup>8</sup>

<sup>6</sup> Assuming that  $N_{tot} \geq \nu$ .

<sup>7</sup> The optimum solution is not always on this line but we found the performance loss from this assumption to be small, as it will be shown in the simulations section.

<sup>8</sup> The assumption  $N_b = \nu$  was termed the *key assumption* in [4] where it was shown to reduce the MMSE-DFE filter computation complexity from cubic to quadratic by using *Cholesky factorization* (instead of inversion) and exploiting matrix structure [5].

For this special case, the error sequence of the FIR MMSE-DFE is *white* when its filter taps are optimized (see the Appendix for a proof).<sup>9</sup>

## IV Numerical Results

We consider the wireless local area network (LAN) environment and compute the loss in  $SNR_{MMSE-DFE,U}$  (as given in (9)) when using the suboptimal 1-dimensional search algorithm instead of the optimal 2-dimensional exhaustive search. The results are shown in Figures 5 and 6 for Channels A and B as specified in the ETSI HIPERLAN standard [6] which model an office environment with maximum delay spread of 390 nsec and 730 nsec, respectively. This corresponds to a channel memory of 8 and 15 for Channels A and B, respectively. The input SNR is set to 15 dB and the results are averaged over 1000 channel realizations. It can be seen that for  $N_{tot} \geq \nu$ , the resulting performance loss is less than 1.0 and 1.4 dB for Channels A and B, respectively. This loss diminishes quickly to less than 0.25 dB for both channels for  $N_{tot} \geq 1.5\nu$ .

## V Conclusions and Future Work

Assuming a fixed total number of taps in the MMSE-DFE, dividing it between feedforward and feedback taps and determining the decision delay that optimize performance require a 2-dimensional exhaustive search. We presented a suboptimal 1-dimensional search algorithm. Simulation results in a wireless LAN environment showed that the proposed reduced-complexity algorithm results in a small performance loss from the optimal solution.

In this paper, we assumed contiguous uniformly-spaced feedforward and feedback taps. For sparse channels, the number of DFE taps can be reduced (without sacrificing performance) by placing them non-uniformly. An interesting open problem is the development of efficient algorithms for determining the optimal number and spacing of the DFE's feedforward and feedback taps when their total number is fixed.<sup>10</sup>

## A Special Case of $N_b = \nu$

A vector of  $N + 1$  channel outputs is given by

$$\mathbf{y}_{i-N}^i = \mathbf{H}_{(N+1) \times (N+\nu+1)} \mathbf{x}_{i-N-\nu}^i + \mathbf{n}_{i-N}^i. \quad (11)$$

We make the following assumptions

<sup>9</sup>Note that this whiteness property was shown for the infinite-length (i.e.  $N = \infty$ ) MMSE-DFE in [2] but it does not hold for arbitrary finite  $(N, N_b, \Delta)$  values.

<sup>10</sup>In the non-uniformly spaced case, this constraint corresponds to a constraint on the computational and implementational complexity, not on the memory or delay requirements.

1. The feedback filter length is equal to the channel memory ( $N_b = \nu$ ). The result also holds for  $N_b > \nu$ , but the additional  $\nu - N_b$  taps are equal to zero.
2. The input data sequence is white ( $\mathbf{R}_{xx} = \mathbf{I}_{N+\nu+1}$ ).
3. The noise sequence is white ( $\mathbf{R}_{nn} = \frac{1}{SNR} \mathbf{I}_{N+1}$ ).
4. Noise and data are uncorrelated.

Under these assumptions, we will show that

1. The error sequence is white.
2. If the decision delay is  $\Delta$ , feedforward filter taps with indices greater than  $\Delta + 1$  are always equal to zero, i.e., the optimum delay is  $\Delta = N$ .

We start with the error symbol  $e_i$  which can be expressed as

$$\begin{aligned}
e_i &= x(i - \Delta) - \underbrace{\begin{bmatrix} w_0^* & \cdots & w_\Delta^* \end{bmatrix}}_{\mathbf{w}_0^{*\Delta}} \underbrace{\begin{bmatrix} y(i) \\ \vdots \\ y(i - \Delta) \end{bmatrix}}_{\mathbf{y}_{i-\Delta}^i} \\
&- \underbrace{\begin{bmatrix} w_{\Delta+1}^* & \cdots & w_N^* \end{bmatrix}}_{\mathbf{w}_{\Delta+1}^{*N}} \underbrace{\begin{bmatrix} y(i - \Delta - 1) \\ \vdots \\ y(i - N) \end{bmatrix}}_{\mathbf{y}_{i-N}^{i-\Delta-1}} \\
&+ \underbrace{\begin{bmatrix} b_1^* & \cdots & b_{N_b}^* \end{bmatrix}}_{\mathbf{b}_1^{*N_b}} \underbrace{\begin{bmatrix} x(i - \Delta - 1) \\ x(i - \Delta - 2) \\ \vdots \\ x(i - \Delta - N_b) \end{bmatrix}}_{\mathbf{x}_{i-\Delta-N_b}^{i-\Delta-1}} \\
&= x(i - \Delta) - \mathbf{w}_0^{*\Delta} (\mathbf{H}_A \mathbf{x}_{i-\Delta-\nu}^i + \mathbf{n}_{i-\Delta}^i) \\
&- \mathbf{w}_{\Delta+1}^{*N} (\mathbf{H}_B \mathbf{x}_{i-N-\nu}^{i-\Delta-1} + \mathbf{n}_{i-N}^{i-\Delta-1}) + \mathbf{b}_1^{*N_b} \mathbf{x}_{i-\Delta-N_b}^{i-\Delta-1} \\
&= x(i - \Delta) \\
&- \mathbf{w}_0^{*\Delta} (\begin{bmatrix} \mathbf{H}_1 & \mathbf{H}_2 \end{bmatrix} \begin{bmatrix} \mathbf{x}_{i-\Delta}^i \\ \mathbf{x}_{i-\Delta-\nu}^{i-\Delta-1} \end{bmatrix} + \mathbf{n}_{i-\Delta}^i) \\
&- \mathbf{w}_{\Delta+1}^{*N} (\begin{bmatrix} \mathbf{H}_3 & \mathbf{H}_4 \end{bmatrix} \begin{bmatrix} \mathbf{x}_{i-\Delta-\nu}^{i-\Delta-1} \\ \mathbf{x}_{i-N-\nu}^{i-\Delta-1} \end{bmatrix} + \mathbf{n}_{i-N}^{i-\Delta-1}) \\
&+ \mathbf{b}_1^{*N_b} \mathbf{x}_{i-\Delta-N_b}^{i-\Delta-1}, \tag{12}
\end{aligned}$$

where  $\mathbf{H}_A$ ,  $\mathbf{H}_B$ , and  $\mathbf{H}_1$ - $\mathbf{H}_4$  are submatrices of  $\mathbf{H}$ . Using the first assumption ( $N_b = \nu$ ), the last equation above becomes

$$\begin{aligned}
e_i &= \underbrace{x(i - \Delta) - \mathbf{w}_0^{*\Delta} (\mathbf{H}_1 \mathbf{x}_{i-\Delta}^i + \mathbf{n}_{i-\Delta}^i)}_{\alpha} \\
&+ \underbrace{(-\mathbf{w}_0^{*\Delta} \mathbf{H}_2 - \mathbf{w}_{\Delta+1}^{*N} \mathbf{H}_3 + \mathbf{b}_1^{*N_b})}_{\beta^*} \mathbf{x}_{i-\Delta-\nu}^{i-\Delta-1} \\
&- \mathbf{w}_{\Delta+1}^{*N} (\mathbf{H}_4 \mathbf{x}_{i-N-\nu}^{i-\Delta-1} + \mathbf{n}_{i-N}^{i-\Delta-1}). \tag{13}
\end{aligned}$$

We want to minimize the MMSE cost  $E[|e_i|^2]$  over all  $\mathbf{w}_0^\Delta$ ,  $\mathbf{w}_{\Delta+1}^N$  and  $\mathbf{b}_1^{N_b}$ . Using the whiteness assumptions on input and noise, the MMSE can be expressed as

$$E[|e_i|^2] = \alpha^2 + \|\beta\|^2 + \mathbf{w}_{\Delta+1}^{*N} (\mathbf{H}_4 \mathbf{H}_4^* + \frac{1}{SNR} \mathbf{I}_{N-\Delta}) \mathbf{w}_{\Delta+1}^N, \tag{14}$$

where  $\alpha$  is a scalar constant and  $\|\cdot\|$  denotes the norm of a vector. It is clear that MMSE attains its minimum value  $\alpha^2$  when

$$\mathbf{w}_{\Delta+1}^N = \mathbf{0} \text{ and } \beta = \mathbf{0} \Rightarrow \mathbf{b}_1^{*N_b} = \mathbf{w}_0^{*\Delta} \mathbf{H}_2. \tag{15}$$

Therefore, the feedforward filter has at most  $\Delta + 1$  nonzero taps, i.e.  $N = \Delta$ . Note that the same conclusion holds had we assumed  $N_b > \nu$  since in that case only the term  $\|\mathbf{b}_{\nu+1}^{*N_b}\|^2$  is added to (14). Hence, to minimize MMSE, we would set  $\mathbf{b}_{\nu+1}^{N_b} = \mathbf{0}$ , i.e., we do not need more than  $\nu$  feedback taps.

From (13) and (15) the error symbol  $e_i$  can be expressed as

$$e_i = x(i - \Delta) - \mathbf{w}_0^{*\Delta} (\mathbf{H}_1 \mathbf{x}_{i-\Delta}^i + \mathbf{n}_{i-\Delta}^i). \tag{16}$$

The optimum  $\mathbf{w}_0^\Delta$  that minimizes (14) can be calculated from (16) as

$$\mathbf{w}_0^{*\Delta} = \mathbf{e}_\Delta^* \mathbf{H}_1^* (\mathbf{H}_1 \mathbf{H}_1^* + \frac{1}{SNR} \mathbf{I}_{\Delta+1})^{-1}, \tag{17}$$

where  $\mathbf{e}_\Delta^*$  is the  $1 \times (\Delta + 1)$  row vector  $[0 \ \cdots \ 0 \ 1]$ . We will now show that  $E[e_j e_i^*] = 0$ . Assume that  $j > i$ . From (16), we have

$$E[e_j e_i^*] = E[e_j (x(i - \Delta) - \mathbf{w}_0^{*\Delta} (\mathbf{H}_1 \mathbf{x}_{i-\Delta}^i + \mathbf{n}_{i-\Delta}^i))^*]. \tag{18}$$

But since  $j > i$ , we get

$$E[e_j x^*(i - \Delta)] = 0. \tag{19}$$

Assume that,  $i - \Delta < j - \Delta \leq i$  (or else the second term in (18) also becomes zero), then

$$\begin{aligned}
E[e_j e_i^*] &= -E[e_j (\mathbf{x}_{i-\Delta}^{*i} \mathbf{H}_1^* + \mathbf{n}_{i-\Delta}^{*i})] \mathbf{w}_0^\Delta \\
&= -E[x(j - \Delta) \mathbf{x}_{i-\Delta}^{*i} \mathbf{H}_1^*] \\
&- \mathbf{w}_0^{*\Delta} (\mathbf{H}_1 \mathbf{x}_{j-\Delta}^j \mathbf{x}_{i-\Delta}^{*i} \mathbf{H}_1^* + \mathbf{n}_{j-\Delta}^j \mathbf{n}_{i-\Delta}^{*i}) \mathbf{w}_0^\Delta. \tag{20}
\end{aligned}$$

The following relations hold

$$\begin{aligned}
E[\mathbf{x}_{j-\Delta}^j \mathbf{x}_{i-\Delta}^{*i}] &= \mathbf{J}; \quad E[\mathbf{n}_{j-\Delta}^j \mathbf{n}_{i-\Delta}^{*i}] = \frac{1}{SNR} \mathbf{J} \\
E[x(j - \Delta) \mathbf{x}_{i-\Delta}^{*i}] &= \mathbf{e}_\Delta^* \mathbf{J}, \tag{21}
\end{aligned}$$

where the  $(\Delta + 1) \times (\Delta + 1)$  matrix  $\mathbf{J}$  has the form  $\mathbf{J} = \begin{bmatrix} \mathbf{0} & \mathbf{0} \\ \mathbf{I}_{\Delta+1-(j-i)} & \mathbf{0} \end{bmatrix}$ . Note that under the assumption  $j - \Delta \leq i$ , we have  $\Delta + 1 - (j - i) > 0$  (or else  $\mathbf{J}$  becomes a zero matrix). Therefore, (20) becomes

$$E[e_j e_i^*] = -(\mathbf{e}_\Delta^* \mathbf{J} \mathbf{H}_1^* - \mathbf{w}_0^{*\Delta} (\mathbf{H}_1 \mathbf{J} \mathbf{H}_1^* + \frac{1}{SNR} \mathbf{J})) \mathbf{w}_0^\Delta. \tag{22}$$

But  $\mathbf{J}\mathbf{H}_1^* = \mathbf{H}_1^*\mathbf{J}$  since  $\mathbf{H}_1^*$  is a square lower-triangular Toeplitz matrix. Therefore,

$$E[e_j e_i^*] = -(\mathbf{e}_\Delta^* \mathbf{H}_1^* - \mathbf{w}_0^{*\Delta} (\mathbf{H}_1 \mathbf{H}_1^* + \frac{1}{SNR} \mathbf{I}_{\Delta+1})) \mathbf{J} \mathbf{w}_0^\Delta. \quad (23)$$

Substituting for  $\mathbf{w}_0^\Delta$  from (17) we get

$$E[e_j e_i^*] = -(\mathbf{e}_\Delta^* \mathbf{H}_1^* - \mathbf{e}_\Delta^* \mathbf{H}_1^*) \mathbf{J} \mathbf{w}_0^\Delta = 0. \quad (24)$$

Similarly, for  $j < i$ ,  $E[e_j e_i^*] = (E[e_i e_j^*])^* = 0$ . Finally, for  $j = i$ , we get

$$\begin{aligned} E[e_i e_i^*] &= E[e_i x^*(i - \Delta)] = 1 - \mathbf{w}_0^{*\Delta} \mathbf{H}_1 \mathbf{e}_\Delta \\ &= 1 - \mathbf{e}_\Delta^* \mathbf{H}_1^* (\mathbf{H}_1 \mathbf{H}_1^* + \frac{1}{SNR} \mathbf{I}_{\Delta+1})^{-1} \mathbf{H}_1 \mathbf{e}_\Delta. \end{aligned}$$

In summary, the error correlation sequence is equal to

$$E[e_i e_j^*] = (1 - \mathbf{e}_\Delta^* \mathbf{H}_1^* (\mathbf{H}_1 \mathbf{H}_1^* + \frac{1}{SNR} \mathbf{I}_{\Delta+1})^{-1} \mathbf{H}_1 \mathbf{e}_\Delta) \delta_{ij},$$

where  $\delta_{ij}$  is the Kronecker delta function (equal to zero for  $i \neq j$  and 1 when  $i = j$ ).

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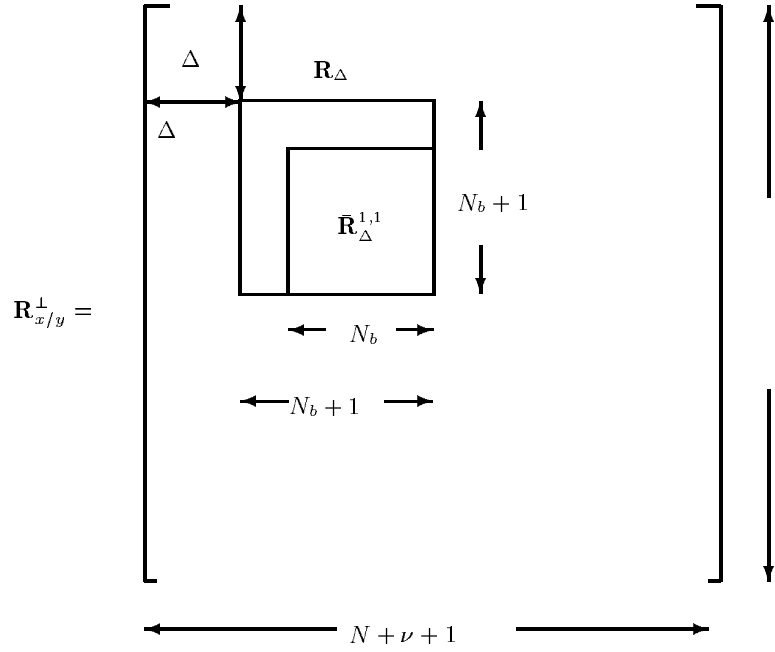


Figure 1: Algebraic interpretation of the parameters  $N$ ,  $N_b$ , and  $\Delta$

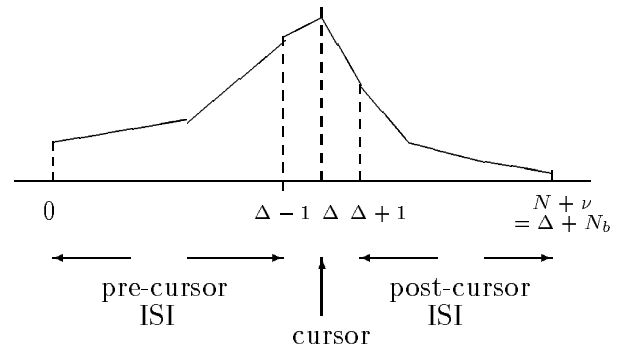


Figure 2: A typical impulse response at the output of the feedforward filter illustrating the assumed relation between the parameters  $N$ ,  $\nu$ ,  $N_b$ , and  $\Delta$

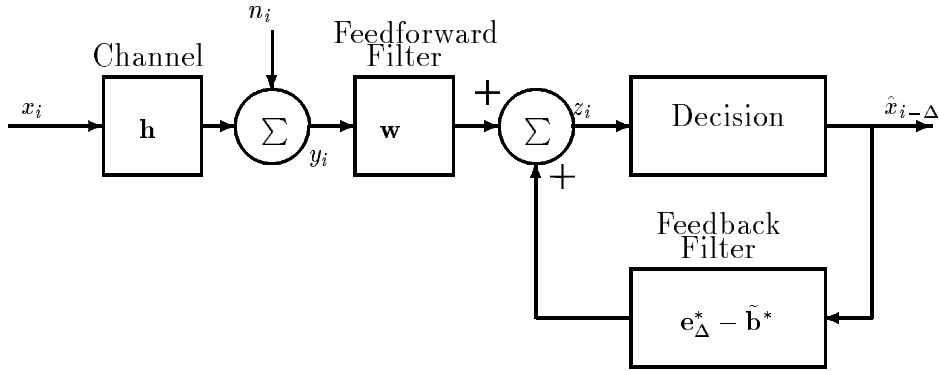


Figure 3: Block diagram of the MMSE-DFE

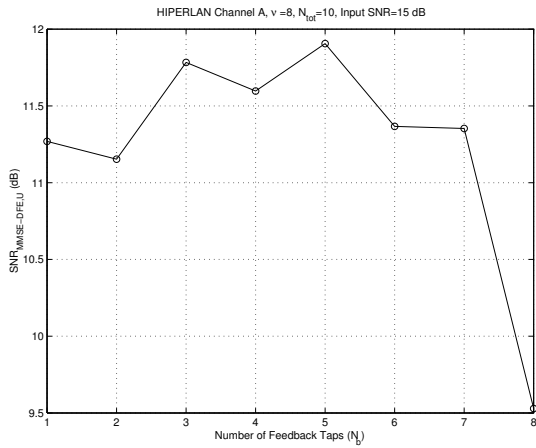


Figure 4: Example of the variation of optimal  $SNR_{MMSE-DFE,U}$  with  $N_b$  for  $N_{tot} = 10$ , input SNR=15 dB on a HIPERLAN A channel

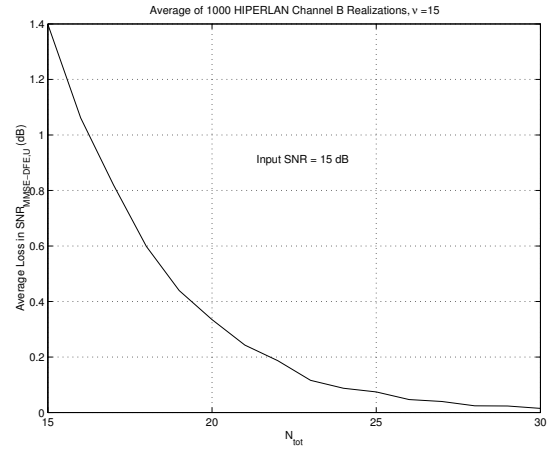


Figure 6: Average loss in  $SNR_{MMSE-DFE,U}$  for the suboptimal algorithm from exhaustive search for the HIPERLAN B channel

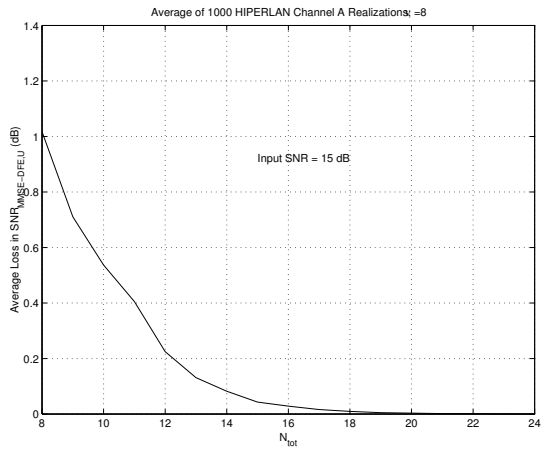


Figure 5: Average loss in  $SNR_{MMSE-DFE,U}$  for the suboptimal algorithm from exhaustive search for the HIPERLAN A channel